

via Riemann-Hurwitz

★★ Plicker formulae

Lecture XXIII

Terminology

$f : S \rightarrow \mathbb{P}^2$ algebraic curve
point coordinates

$f^* : S \rightarrow \mathbb{P}^{2*}$ dual curve
line coordinates
envelope of tangents

$(f^*)^* = f$

• regular point:

smooth point of f and f^*

local expression

$f(z) = [1, z + \dots, z^2 + \dots]$

x_0, x_1, x_2
 $p: [1, 0, 0]$

• ordinary flex

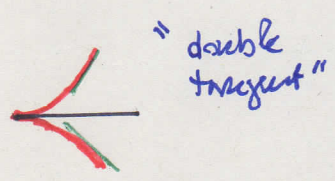
smooth point whose tangent line has contact of order 3

$f(z) = [1, z + \dots, z^3 + \dots]$

$f^*(z) = [1, z^2 + \dots, z^3 + \dots] \rightarrow$ cusp of f^*

• cusp of f

$f(z) = [1, z^2 + \dots, z^3 + \dots]$

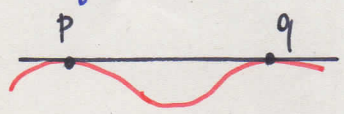


p is a flex of $f \iff p$ is a cusp for f^*

• bitangent of f

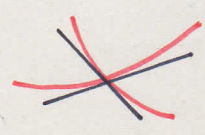
p , not a flex, such that its tangent simply touches f in another point $q \neq p$

tangent simply touches f in another point $q \neq p$



• ordinary double point (node)

two branches of f cross transversally



p is a bitangent for $f \iff p$ is an ordinary double point for f^*

Notation

g : genus of S

$C = f(S) \quad C^* = f^*(S)$

d, d^* : degrees of C, C^* resp.

d^* : class of C

b, b^* : # bitangents (of C, C^*)

(# of tangents emanating from a generic pt)

f, f^* : # flexes (of C, C^*)

r, r^* : # cusps (of C, C^*)

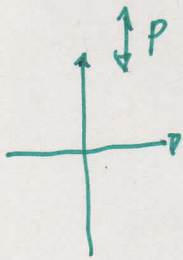
δ, δ^* : # double points of C, C^*

Clearly:

$$\left[\begin{array}{l} b = \delta^* \quad , \quad b^* = \delta \\ f = r^* \quad , \quad f^* = r \end{array} \right]$$

Assumptions: p generic, not lying on tangents issuing from the singular points

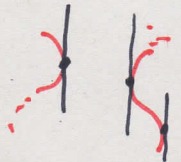
local coordinates $[x_0, x_1, x_2]$ on \mathbb{P}^2 , $p: [0, 0, 1] = Y_0$



$C: g(x_0, x_1, x_2) = 0 \quad \deg C = d$

tangent lines to C from p \leftrightarrow smooth

points of C such that $\frac{\partial g}{\partial x_2}(q) = 0$



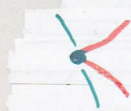
Let $C': \frac{\partial g}{\partial x_2} = 0 \quad \deg C' = d-1$

passes through each bitangent of C^* & flex of C^* & cusp of C with intersection multiplicity 2 & 3 resp.

at a sing. pt: $\frac{\partial g}{\partial x_i} = 0$



2 intersections



3 intersections

Then
 # points of intersection of C' with $C \setminus \{\text{singular pts}\}$
 $= (C \cdot C') - 2\delta - 3k$, namely

$$\boxed{d^* = d(d-1) - 2\delta - 3k}$$

class of G
 $=$ # tangents
 emanating
 from a generic
 pt
 in our case p

Subsequently, project
 C onto a line \mathbb{P}^1 :

$$\pi : C \rightarrow \mathbb{P}^1$$

d -sheeted covering of \mathbb{P}^1

Therefore, by Riemann-Hurwitz

$$\chi(S) = 2 - 2g = 2d - b$$

branch points of
 π of $S \rightarrow \mathbb{P}^1$

Now, a smooth pt of G , say q

is a branch pt iff $\frac{\partial g}{\partial x_2}(q) = 0$



\Rightarrow among the smooth pts of G one has

$$d(d-1) - 2\delta - 3k \text{ branch pts}$$

Now, an ordinary double point does not yield a
branch pt, and a cusp is a branch pt of order 1

important!

(3)

XXIII - 3

hence

$$\begin{aligned}
 b &= d(d-1) - 2\delta - 3r + r \\
 \# \text{ branch } \# &= d(d-1) - 2\delta - 2r
 \end{aligned}$$

Thus

$$\begin{aligned}
 2 - 2g &= 2d - \overbrace{d(d-1)}^{-b} + 2\delta + 2r \\
 &= 2d - d^2 + d + 2\delta + 2r \\
 &= 3d - d^2 + 2\delta + 2r
 \end{aligned}$$

$$\begin{aligned}
 \Rightarrow 2g &= 2 - 3d + \overbrace{d^2}^{(d-1)(d-2)} - 2\delta - 2r = (d-1)(d-2) - 2\delta - 2r \\
 g &= \frac{(d-1)(d-2)}{2} - \delta - r
 \end{aligned}$$

bitangents $b = \delta^*$
 flexes $f = r^*$

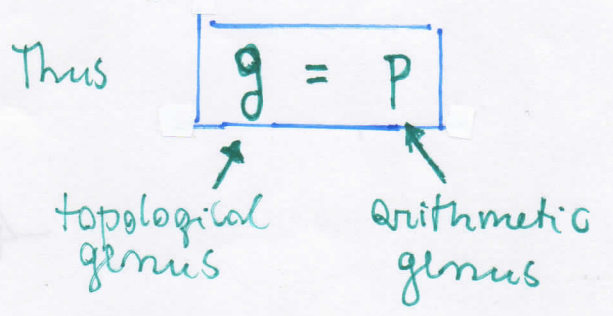
$$d^* = d(d-1) - 2\delta - 3r \qquad g = \frac{(d-1)(d-2)}{2} - \delta - r$$

Dually

$$d = d^*(d^*-1) - 2b - 3f \qquad g^* = g = \frac{(d^*-1)(d^*-2)}{2} - \delta^* - r^2$$

$$= \frac{(d^*-1)(d^*-2)}{2} - b - f$$

★ ★ Plücker formulae



Lüroth's Theorem (classical proof)

Let $\mathcal{L} : x_i = \frac{\phi_i(t)}{\phi_0(t)} \quad i=1 \dots n$ (\mathcal{L} rational)

If every point of \mathcal{L} arises from $\mathbb{R} \geq 1$ values of t ,
then $\exists \tau = \tau(t)$ (rational function) such that
the correspondence $\tau \leftrightarrow$ pts of \mathcal{L} is 1-1.

Proof. Let the ϕ_i without common (non constant) factor
and with the same degree. Let t_1 a generic value
of t , and $P_{t_1} = \left(\frac{\phi_1(t_1)}{\phi_0(t_1)}, \dots, \frac{\phi_n(t_1)}{\phi_0(t_1)} \right)$

The values $t_2 \dots t_k$ yielding the same P_{t_1} are such that

$$\text{rank} \begin{pmatrix} \phi_1(t) & \dots & \phi_n(t) \\ \phi_1(t_1) & \dots & \phi_n(t_1) \end{pmatrix} = 1, \text{ namely, the common}$$

zeros of $\Delta_{ij}(t) := \phi_i(t_1)\phi_j(t) - \phi_i(t)\phi_j(t_1) \quad \begin{matrix} i, j=1 \dots n \\ i \neq j \end{matrix}$

i.e. the zeros of their GCD

$$\Delta(t, t_1) := a_0(t_1)t^n + a_1(t_1)t^{n-1} + \dots + a_k(t_1)$$

Upon varying t_2 , the coefficients are not constant, otherwise
the roots of Δ would be independent of t_2 , and this
is absurd since t_1 is a root.

We can also assume that all roots of Δ vary with t_2
(otherwise, we could remove those staying constant).

Therefore, for generic t_1 , $\Delta(t, t_2)$ has degree n in t
(and $a_0(t_1) \neq 0$). Any other value $t_2 \dots t_n$ yielding
the same point would have worked equally well
 $\tau(t_1)$

Hence, there exists j ($j=1$, say)

such that $\frac{a_j(t_1)}{a_0(t_1)}$ depends effectively on t_2

and
$$\frac{a_1(t_1)}{a_0(t_1)} = \frac{a_2(t_2)}{a_0(t_2)} = \dots = \frac{a_k(t_k)}{a_0(t_k)}$$

Then $\tau := \frac{a_1(t)}{a_0(t)}$ is the sought new parameter

From
$$\begin{cases} \tau a_0(t) - a_1(t) = 0 \\ \alpha_i \phi_0(t) - \phi_i(t) = 0 \end{cases}$$

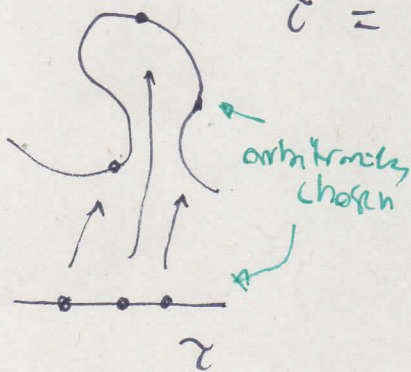
one has, after eliminating t ,

$$g(\alpha_i, \tau) = 0 \quad (\text{algebraic})$$

For every τ , α_i is the i -th coordinate of the point $P(\tau)$

All acceptable τ 's are obtained via

$$\tau' = \frac{a\tau + b}{c\tau + d} \quad ad - bc \neq 0$$



Other definition:

Rational curve in \mathbb{P}^r :

$$\alpha_i = f_i(\lambda) \quad i = 0 \dots r$$

polynomials without a common factor

* A rational curve is algebraic: for simplicity let $r=2$ (plane curve); one has

$$\begin{aligned} f_1(\lambda)\alpha_0 - f_0(\lambda)\alpha_1 &= 0 \\ f_2(\lambda)\alpha_0 - f_0(\lambda)\alpha_2 &= 0 \end{aligned}$$

Elimination of λ leads to a homogeneous equation $F(\alpha_0, \alpha_1, \alpha_2) = 0$

★ Luröth's theorem

"topological" version

compact RS

Let $\pi: \mathbb{P}^1 \rightarrow S$ be a d -sheeted branched covering. Then $S \approx \mathbb{P}^1$.

Proof. From Riemann-Roch we have:

$$\chi(\mathbb{P}^1) = d \cdot \chi(S) - b$$

" 2
" 2

$$\Rightarrow d \cdot \chi(S) = 2 + b$$

$$\Rightarrow \chi(S) > 0$$

$$\Rightarrow g = 0$$

Thus S is \mathbb{P}^1 .